

# SECURE, EFFICIENT AND FLEXIBLE DATA SHARING ON DISTRIBUTED CLOUDS USING KEY AGGREGATION

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*Abstract— Cloud computing technology is wide used so the info are often outsourced on cloud will accessed simply. totally different members can share that knowledge through totally different virtual machines however gift on single physical machine. however the factor is user don't have physical management over the outsourced knowledge. the requirement is to share knowledge firmly among users. The cloud service supplier and users authentication is important to create positive no loss or leak of users knowledge. Privacy conserving in cloud is very important certify the users identity isn't unconcealed to everybody. On cloud anyone will share knowledge the maximum amount they need to i.e. solely selected content are often shared. Cryptography helps the info owner to share the info to in safe method. thus user encrypts knowledge and uploads on server. Different secret writing and decipherment keys square measure generated for various knowledge. The secret writing and decipherment keys is also totally different for different set of information. solely those set of decipherment keys square measure shared that the chosen knowledge are often decrypted. Here a public-key cryptosystems that generate a ciphertext that is of constant size. so to transfer the decipherment rules for range of ciphertext. The distinction is one will collect a group of secret keys and create them as little size as one key with holding a similar ability of all the keys that square measure fashioned in an exceedingly cluster. This compact mixture key are often with efficiency sent to others or to be keep in an exceedingly smart card with very little secure storage.*

*Key words: computing, key aggregate, information, owner, Keys, knowledge.*

## I. INTRODUCTION

Cloud computing is wide increasing technology; information will be saved on cloud remotely and may have access to large applications with quality services that are shared among customers. As increase in outsourcing of

information the cloud computing serves will the management of information. Its versatile and value optimizing characteristic motivates the end user moreover as enterprises to store the info on cloud. The business executive attack is one in all security concern which's wants to be centered. Cloud Service supplier ought to ensure whether audits are command for users World Health Organization have physical access to the server. As cloud service supplier stores the info of different users on same server it's potential that user's private information is leaked to others. the general public auditing system of data storage security in cloud computing provides a

privacy-preserving auditing protocol. It is necessary to create certain that the info integrity while not compromising the namelessness of the info user. to confirm the integrity the user will verify information on their information, upload and verify information. The main concern is a way to share the info firmly the answer is cryptography. The question is however will the encrypted information is to be shared. The user should offer the access rights to the opposite user because the information is encrypted and the decoding key ought to be send firmly. For AN example Alice keeps her personal information i.e. photos on drop box and she or he doesn't need to share it with everybody. because the assailant could access the info therefore it's inconceivable to deem predefine

privacy protective mechanism therefore she all the photos were encrypted by her on secret writing key whereas uploading it.

Suppose some day she needs to share few photos along with her friend Bob, either she will be able to write all photos with one key and send to him or she will be able to produce

write with totally different keys and send it. The un-chosen knowledge is also leaked to Bob if the single key generated for secret writing therefore produce distinct keys of data and send single key for sharing. A new manner for public-key secret writing is employed as known as keyaggregate cryptosystem (KAC).

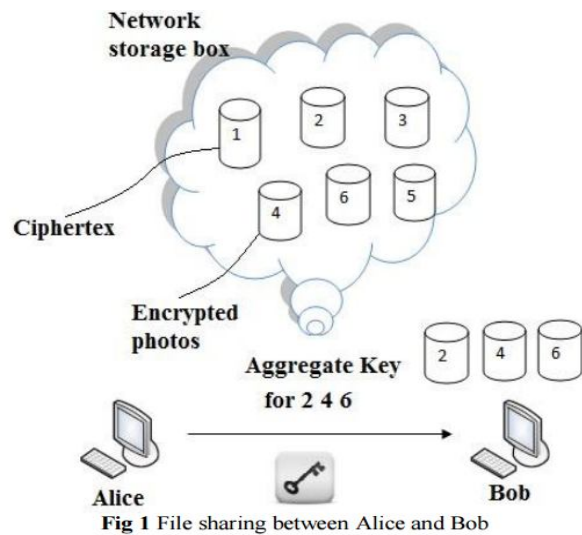


Fig 1 File sharing between Alice and Bob

The secret writing is completed through Associate in Nursing symbol of Ciphertext called category, with public key. The categories area unit fashioned by classifying the ciphertext. The key owner has the master secret key that is helpful for extracting secret key. therefore in higher than senario currently the aice will send a combination key to bob through a email and the encrypted knowledge is downloaded from dropbox through the aggregate key.This is shown in figure1.

## II. RELATED WORK

### SYMMETRIC-KEY ENCRYPTION WITH COMPACT KEY:

Benaloh et al presented Associate in Nursing cryptography haem that's originally projected for briefly causing sizable quantity of keys in broadcast scenario. the event is simple which we tend to briefly review its key derivation methodology here for a concrete description of what square measure the fascinating properties we tend to would like to achieve. The derivation of the key for a group of classes (which can be a group of all achievable ciphertext classes) is as follows. A composite modulus is chosen where  $p$  and  $q$  are two large random primes. A master

secret secret is chosen willy-nilly. each class is expounded to a particular prime. of those prime numbers is place inside the general public system parameter. A constant-size key for set is generated. For those who area unit delegated the access rights for  $S'$  is generated. However, it's designed for the symmetric-key setting instead. The content provider should get the corresponding secret keys to cypher info, that is not applicable for many applications. as a results of technique is used to come back up with a secret value rather than a strive of public/secret keys, it's unclear the thanks to apply this idea for public-key cryptography theme. Finally, we tend to tend to notice that there square measure schemes that try to reduce the key size for achieving authentication in symmetric-key cryptography. However, sharing of decryption power is n't a priority in these schemes.

### a. IBE WITH COMPACT KEY

Identity-based cryptography (IBE) ( could also be a public-key cryptography throughout that the public-key of a user are set as Associate in Nursing identity-string of the user (e.g., Associate in Nursing email address, mobile number). there is a personal key generator (PKG) in IBE that holds a master-secret key and issues a secret key to each user with connexion the user identity. The content provider can take the overall public parameter and a user identity to put in writing in code a message. The recipient can rewrite this ciphertext by his secret key. Guo et al, tried to create IBE with key aggregation. In their schemes, key aggregation is strained inside the sense that every one keys to be mixture ought to return from fully totally different —identity divisionsl. whereas there area unit Associate in Nursing exponential form of identities then secret keys, exclusively a polynomial variety of them are mixture.This significantly can increase the costs of storing and causing ciphertexts, that is impractical in many things like shared cloud storage. As in our own thanks to attempt to to the current is commonly to use hash perform to the string denoting the category, and keep hashing repeatedly until a primary is obtained as a result of the output of the hash perform. we tend to tend to mentioned, our schemes feature constant ciphertext size, and their security holds inside the conventional model. In fuzzy

IBE, one single compact secret key can rewrite ciphertexts encrypted below many identities that area unit go a specific topological space, but not for Associate in Nursing arbitrary set of identities and so it does not match with our set up of key aggregation.

**b. ATTRIBUTE-BASED ENCRYPTION**

Attribute-based coding (ABE) permits every ciphertext to be related to associate degree attribute, and also the aster-secret key holder can extract a secret key for a policy of those attributes in order that a ciphertext are often decrypted by this key if its associated attribute conforms to the policy. as an example, with the key key for the policy(1 ∨ 3 ∨ 6 ∨ 8) , one will decipher ciphertext labeled with category 1, 3, 6 or 8. However, the main concern in ABE is collusion-resistance however not the compactness of secret keys. Indeed, the scale of the key typically will increase linearly with the quantity of attributes it encompasses, or the cipher text-size isn't constant.

| Different Schemes                         | Ciphertext size | Decryption key size | Encryption type         |
|---|-----------------|---------------------|-------------------------|
| Key assignment schemes                    | Constant        | Non-constant        | Symmetric or public-key |
| Symmetric-key encryption with compact key | Constant        | Constant            | Symmetric key           |
| IBE with compact key                      | Non-constant    | Constant            | Public key              |
| Attribute based encryption                | Constant        | Non-constant        | Public key              |
| KAC                                       | Constant        | Constant            | Public key              |

Comparison between KAC scheme and other related scheme

**KEY-AGGREGATE CRYPTOSYSTEM**

In key-aggregate cryptosystem (KAC), users encipher a message not solely beneath a public-key, however conjointly beneath associate symbol of ciphertext referred to as category. which means the ciphertexts area unit additional categorised into totally different categories. The key owner holds a master-secret called master-secret key, which may be accustomed extract secret keys for various categories. a lot of significantly, the extracted key have will be associate

combination key that is as compact as a secret key for one category, however aggregates the ability of the many such keys, i.e., the decryption power for any set of ciphertext categories.

With our example, Alice will send Bob one combination key through a secure e-mail. Bob will transfer the encrypted photos from Alice's Box.com area so use this combination key to rewrite these encrypted knowledge. The sizes of ciphertext, public-key, master-secret key and combination key in KAC schemes area unit all of constant size. the general public system parameter has size linear within the variety of ciphertext categories, but solely atiny low a part of it's required anytime and it will be fetched on demand from massive (but non-confidential) cloud storage.

**III. FRAME WORK**

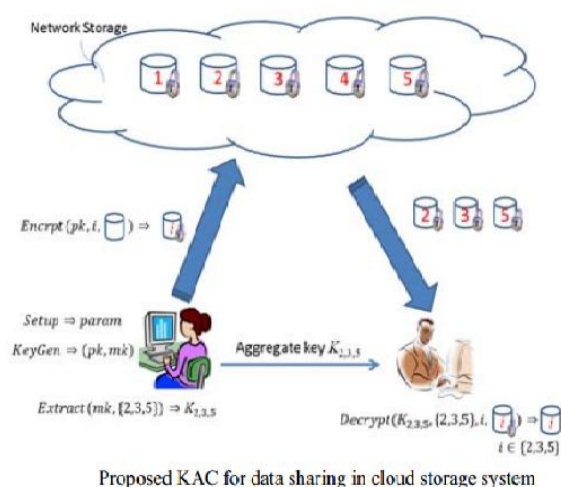
The data owner establishes the general public system parameter through Setup and generates a public/master-secret key try through KeyGen. information are often encrypted via write by anyone United Nations agency conjointly decides what ciphertext category is related to the plaintext message to be encrypted. the info owner will use the master-secret key try to get AN mixture coding key for a group of cipher text classes through Extract. The generated keys are often passed to delegates firmly through secure e-mails or secure devices Finally, any user with AN mixture key will rewrite any cipher text only if the cipher text's category is contained within the mixture key via Decrypt. Key mixture coding schemes accommodates 5 polynomial time algorithms as follows:

- 1.Setup (1 λ , n) : The data owner establish public system parameter via Setup. On input of a security level parameter 1 λ and number of ciphertext classes n , it outputs the public system parameter *param*.
2. KeyGen: It is executed by data owner to randomly generate a public/ master-secret key pair (Pk, msk).
3. Encrypt (pk, i, m) : It is executed by data owner and for message m and index i ,it computes the ciphertext as C.
4. Extract (msk, S): It is executed by data owner for delegating the decrypting power for a certain set of ciphertext classes and it outputs the aggregate key for set S denoted by Ks.

5. Decrypt ( $K_s, S, I, C$ ): It is executed by a delegate who received, an aggregate key  $K_s$  generated by Extract. On input  $K_s$ , set  $S$ , an index  $i$  denoting the ciphertext class ciphertext  $C$  belongs to and output is decrypted result  $m$ .

### SHARING ENCRYPTED DATA

A canonical application of KAC is knowledge sharing. The key aggregation property is very helpful after we expect delegation to be efficient and versatile. The KAC schemes change a content supplier to share her knowledge in a very confidential and selective approach, with a hard and fast and small ciphertext growth, by distributing to every approved user one and little mixture key.



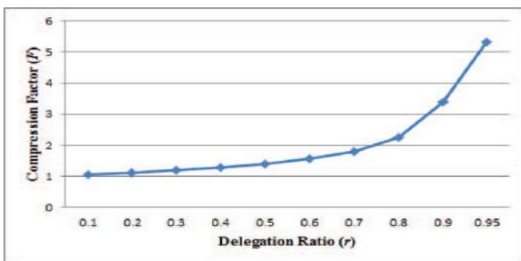
Data sharing in cloud storage victimisation KAC, illustrated in Figure one . Suppose Alice desires to share her knowledge money supply , $m_2, \dots, m_n$  on the server. She 1st performs Setup ( $1 \lambda, n$ ) to induce param and execute KeyGen to induce the public/master-secret key combine ( $pk, msk$ ). The system parameter param and public-key  $pk$  are often created public and master-secret key  $msk$  ought to be unbroken secret by Alice. Anyone will then encrypt every  $m_i$  by  $C_i = \text{cipher}(pk, i, m_i)$ . The encrypted knowledge area unit uploaded to the server. With param and  $pk$ , folks that join forces with Alice will update Alice's knowledge on the server. Once Alice is willing to share a group  $S$  of her knowledge with a devotee Bob, she will be able to calculate the aggregate key  $K_{s, S}$  for Bob by acting Extract ( $msk, S$ ). Since  $K_{s, S}$  is simply a continuing size key, it's simple to be sent to Bob through a secure e-mail. once getting the combination key, Bob will transfer the information he's

approved to access. That is, for every  $i \in S$ , Bob downloads  $C_i$  from the server. With the combination key  $K_{s, S}$ , Bob will decipher every  $C_i$  by decipher ( $K_{s, S}, i, C_i$ ) for every  $i \in S$ .

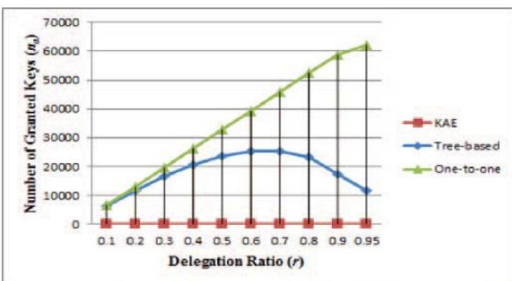
### IV. EXPERIMENTAL RESULTS

Our methodologies change the pressure issue ( $F = n$  in our plans) to be a tunable parameter, at the expense of  $O(n)$ -estimated framework parameter. cryptography is tired consistent time, while coding is tired  $O(|S|)$  bunch augmentations (or reason expansion on elliptic bends) with 2 blending operations, where  $S$  is that the situated of ciphertext classes decryptable by the allowed blend key and  $|S| \leq n$ . obviously, key extraction needs  $O(|S|)$  group increases also, that a substitution progress on the stratified key task (an old approach) that jam regions giving the wholes of the key-holders offer comparative edges is our methodology of "compacting" mystery keys openly key cryptosystems. These open key cryptosystems production figure writings of consistent size ostensible efficient designation of mystery composing rights for any arrangement of figure writings is conceivable. This not solely upgrades client security and classifiedness of information in cloud capacity, however it'll this by supporting the dissemination ornaming of mystery keys changed for diverse figure content classes and producing keys by various determination of figure content class properties of the data and its related keys. This aggregates up the extent of our paper. As there is an utmost assault choice the amount the quantity of figure content classes already & notwithstanding the exponential development inside the amount of figure messages in distributed storage, there is an interest for reservation of figure content classes for future utilization. With respect to potential alterations furthermore, upgrades to our present reason, in future, the parameter size region unit generally modified ostensible its independent the very pinnacle of style of figure content classes. to boot, a uniquely composed cryptosystem, with the livelihood of a precise security equation, as partner degree illustration, the Diffie-Hellman Key-Exchange strategy, which can at that point be imperviable, or at the principal confirmation against

overflowing at the part of prudent key naming, will affirm that one can transport same keys on cell phones without apprehension of overflow.



Compression achieved by the tree-based approach for delegating different ratio of the classes



Number of granted keys ( $n_a$ ) required for different approaches in the case of 65536 classes of data

## V. CONCLUSION

Another progress on the class-cognizant key task (a antiquated methodology) that jelly regions giving the sums of the key-holders offer comparative edges is our methodology of "packing" mystery keys freely key cryptosystems. These open key cryptosystems assembling figure writings of consistent size indicated sparing assignment of mystery composing rights for any arrangement of figure writings is practical. This not exclusively upgrades client security and secrecy of learning in distributed storage, on the other hand it will this by supporting the circulation or delegating of mystery keys various for diverse figure content classifications and creating keys by different deduction of figure content classification properties of the information and its related keys. This wholes up the extent of our paper. As there's a point of confinement assault assortment the sum the quantity of figure content classes previously & not to mention the exponential development inside the quantity of figure messages in distributed storage, there's a necessity for reservation of figure content classes for future utilization. With respect to potential alterations and improvements to our current reason, in future, the parameter size are regularly

adjusted indicated its independent the very pinnacle of mixture of figure content classifications. to boot, an extraordinarily composed cryptosystem, with the usage of a right security algorithmic guideline, as a sample, the Diffie-Hellman Key-Exchange philosophy, which might at that point be step confirmation, or at the most evidence against overflowing along the edge of practical key delegating, can verify that one will transport same keys on portable gadgets without stressing of overflowing.

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